

| Side-Channel Security

Chapter 7: Network Side Channels

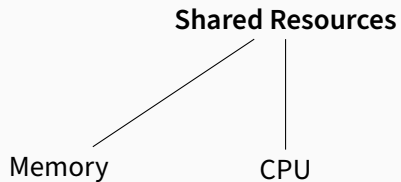
Stefan Gast

2025-04-03

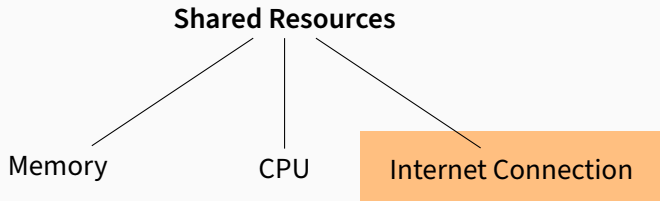
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Introduction

What to Attack?



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Demo: Network Traffic Depends on Activity

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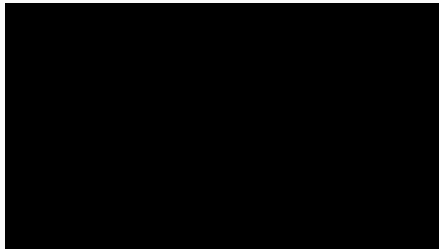
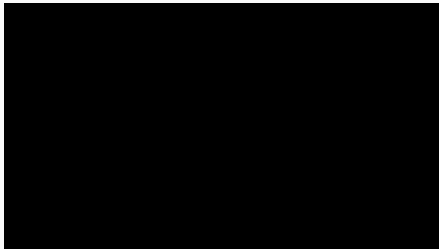
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 - ...
- attacker-in-the-middle, mostly used against privacy-enhancing tunnels

Which video segment uses more bandwidth?



<https://www.youtube.com/watch?v=LNI8rnxxVvQ>

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- fingerprint!

- Reed and Kranch, 2017 [RK17]: Netflix
- Schuster et al. , 2017 [SST17]: YouTube, Netflix, Amazon, Vimeo
- Gu et al. , 2018 [Gu+18]: self-hosted DASH server
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 - ...
- attacker-in-the-middle or with JavaScript

- SSH keystroke timings [SWT01]

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- deanonymization of Tor users [RSG98; AYR15; Wan+11]
- language [Wri+07] and phonemes [Whi+11] of VoIP calls
- other privacy-critical information [Che+10; LM18]

SnailLoad: Remote Traffic Analysis via TCP [Gas+24]

**Some of you probably know the
effect...**

- DSL, Fiber, LTE, 5G: different throughput

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- buffering before last mile is necessary!

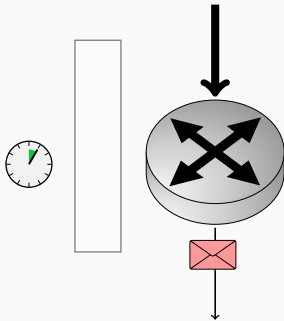


Figure 1: Connection idle

Packet Buffering

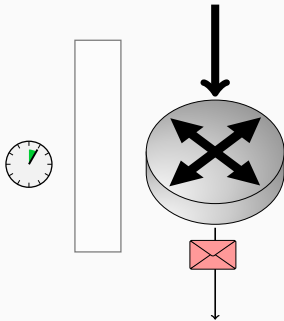


Figure 1: Connection idle

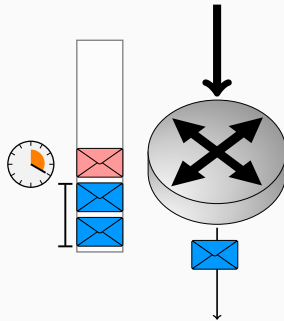


Figure 2: Connection busy

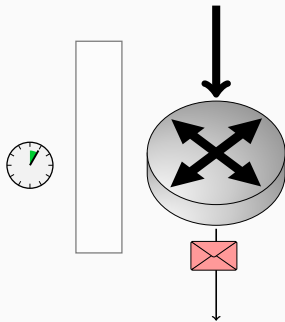


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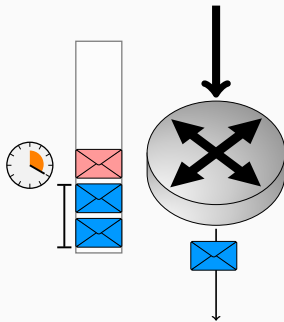


Figure 2: Connection busy

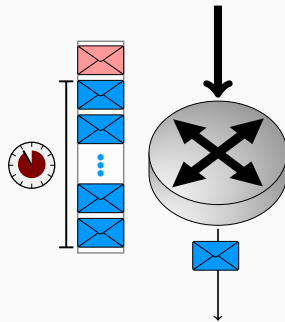


Figure 3: Bufferbloat

Network Activity Causes Latency Spikes

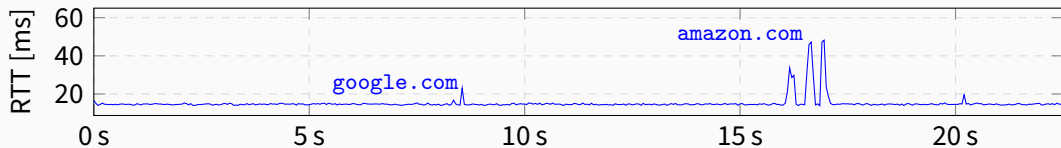


Figure 4: Same machine pinging 8.8.8.8

Network Activity Causes Latency Spikes

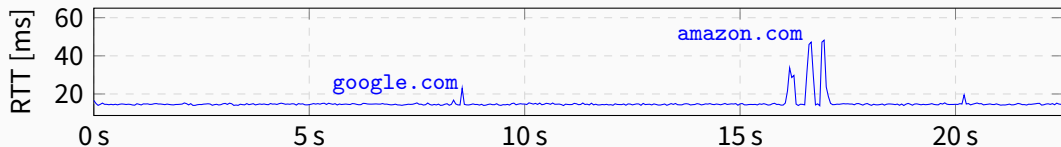


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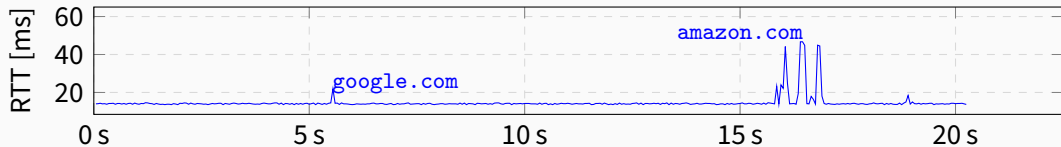


Figure 5: Different machine sharing the same internet connection pinging 8.8.8.8

Idle and Busy Round-Trip-Times

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Idle and Busy Round-Trip-Times

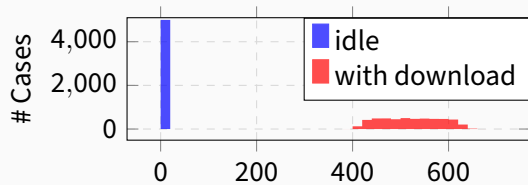


Figure 6: RTT [ms], ADSL-1, 50 Mbit/s

Idle and Busy Round-Trip-Times

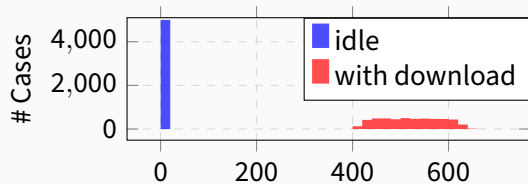


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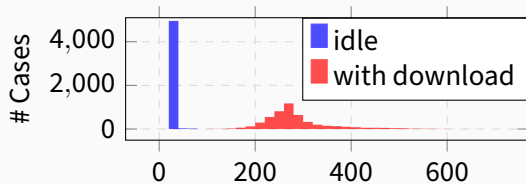


Figure 7: RTT [ms], LTE, 75 Mbit/s

Idle and Busy Round-Trip-Times

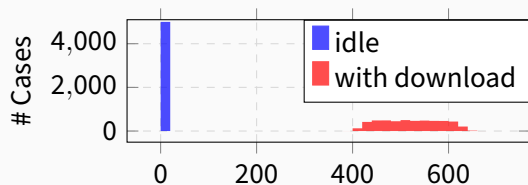


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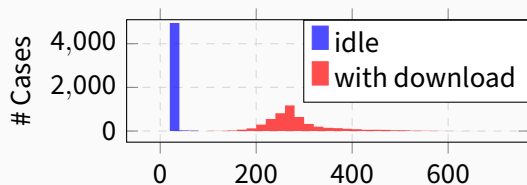


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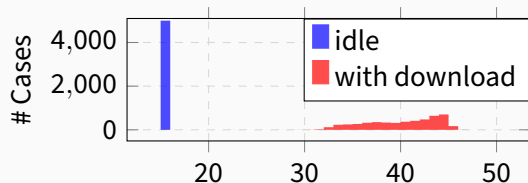


Figure 8: RTT [ms], FTTH-1, 80 Mbit/s

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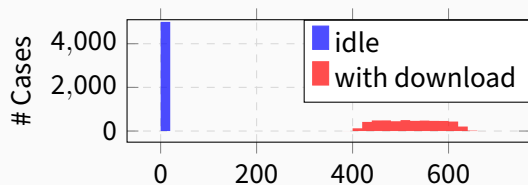


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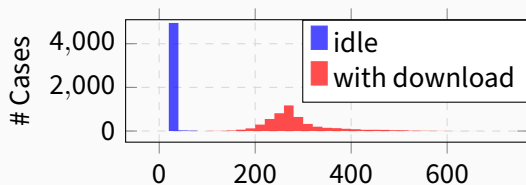


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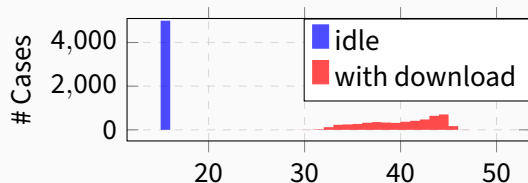


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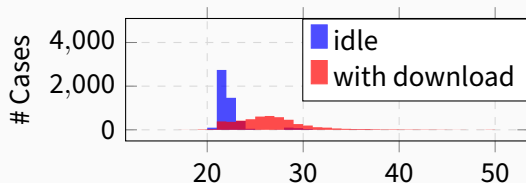
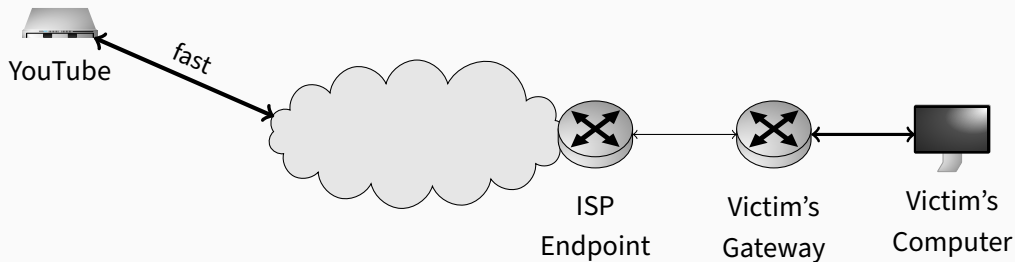
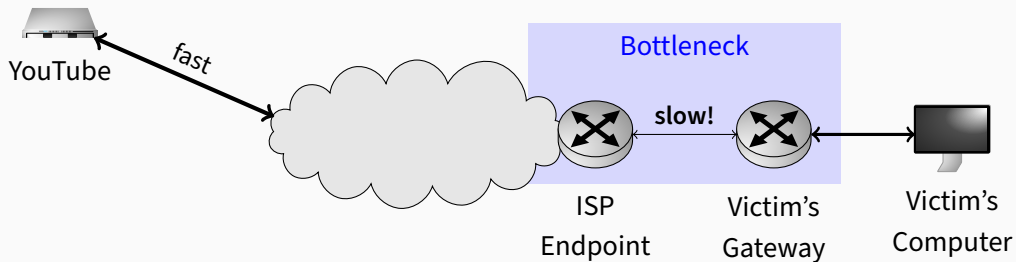


Figure 9: RTT [ms], Cable, 80 Mbit/s

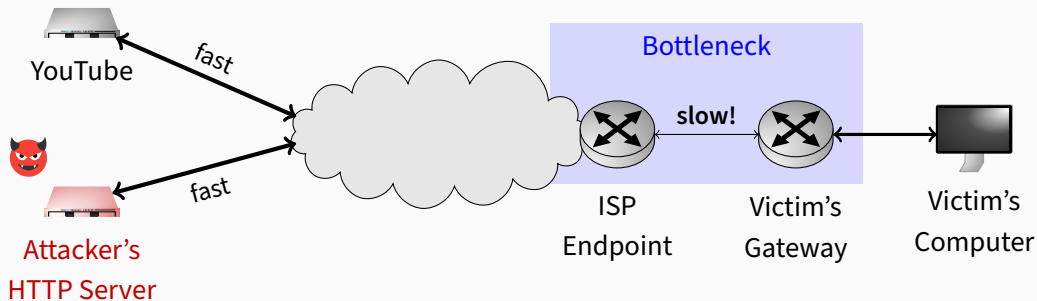
Attack Setup



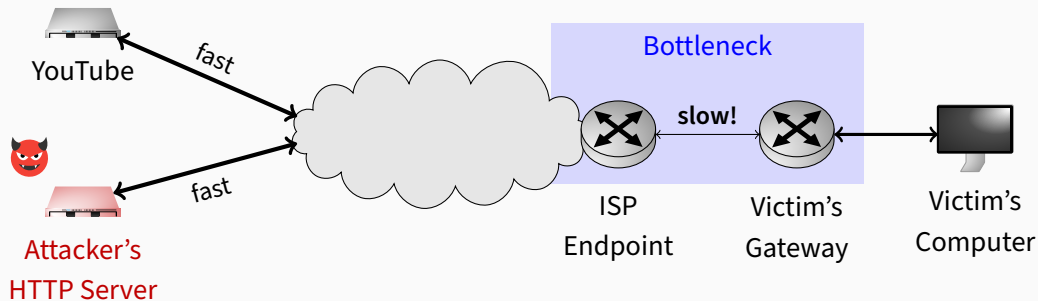
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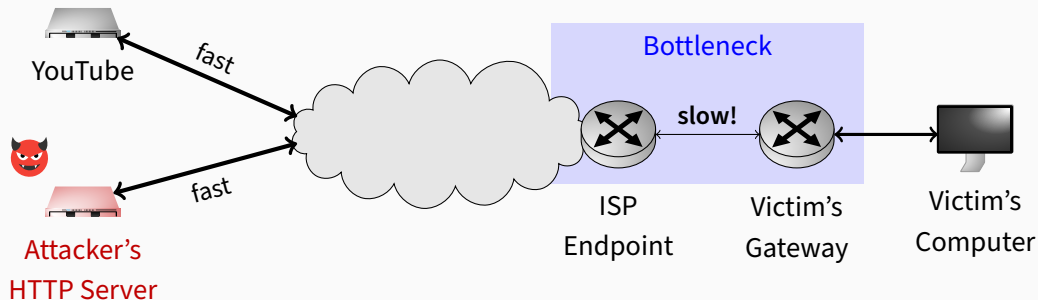


Attack Setup



- Various scenarios: Compromised websites, malicious ads, emails, and more

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- Different ways attackers can exploit network traffic to perform attacks

Polling the Server's Send Buffer To Measure RTTs

begin

acked \leftarrow **false**;

start \leftarrow get_current_time();

send(sock, b, 1, 0);

repeat

if ioctl(sock, SIOCOUTQ) = 0 **then**

 acked \leftarrow **true**;

end

until acked;

end \leftarrow get_current_time();

return end - start;

end

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Polling the Server's Send Buffer To Measure RTTs

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        |     acknowledged ← true;  
        | end  
    until acknowledged;  
    end ← get_current_time();  
    return end - start;
```

end

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- pre-process traces with an STFT
- KERAS (Tensorflow)

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Table 1: CNN Parameters

Type	Parameters	Activation
Conv2D	filters=32, kernel size=[5,5], strides=[1,1]	ReLU
MaxPooling2D	pool size=[2,2], strides=[2,2]	-
Conv2D	filters=64, kernel size=[3,3], strides=[1,1]	ReLU
MaxPooling2D	pool size=[2,2], strides=[2,2]	-
Conv2D	filters=128, kernel size=[3,3], strides=[1,1]	ReLU
MaxPooling2D	pool size=[2,2], strides=[2,2]	-
Flatten	-	-
Dense	output size=1024	ReLU
Dense	output size=512	ReLU
Dense	output size=10	Softmax

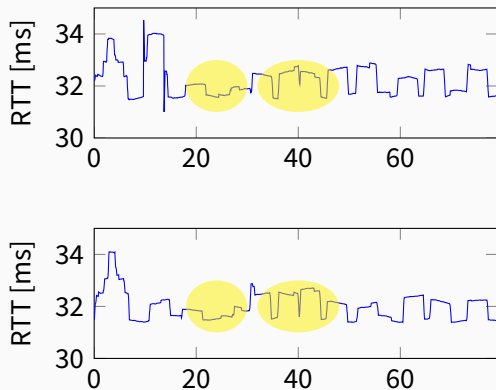


Figure 10: Video A, Time in seconds on x axis

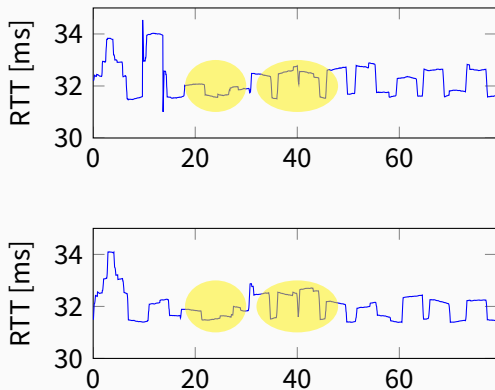


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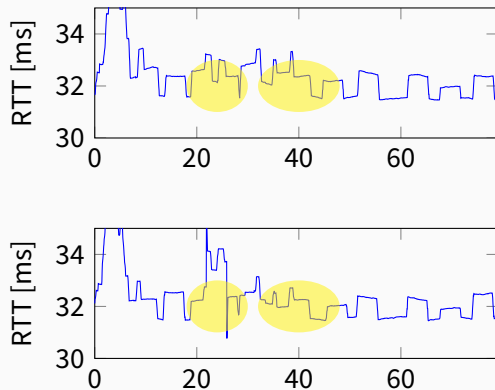
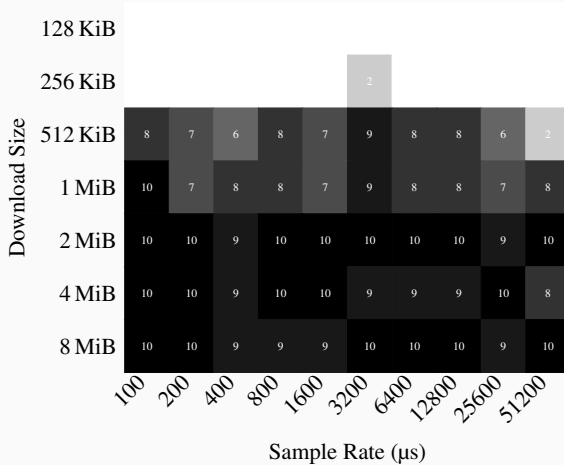
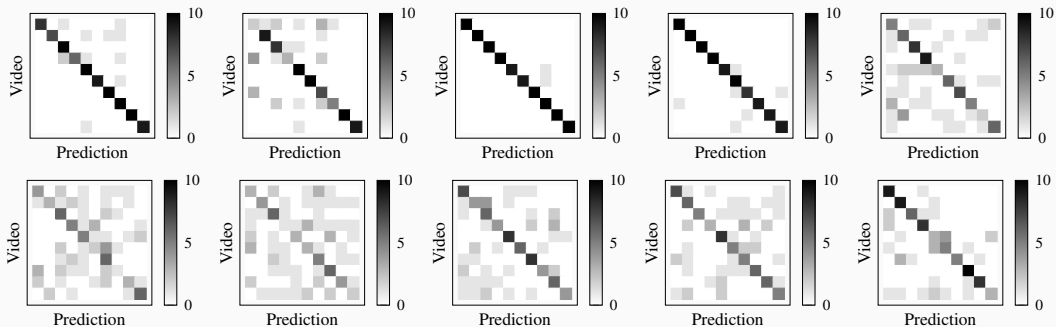


Figure 11: Video B, Time in seconds on x axis

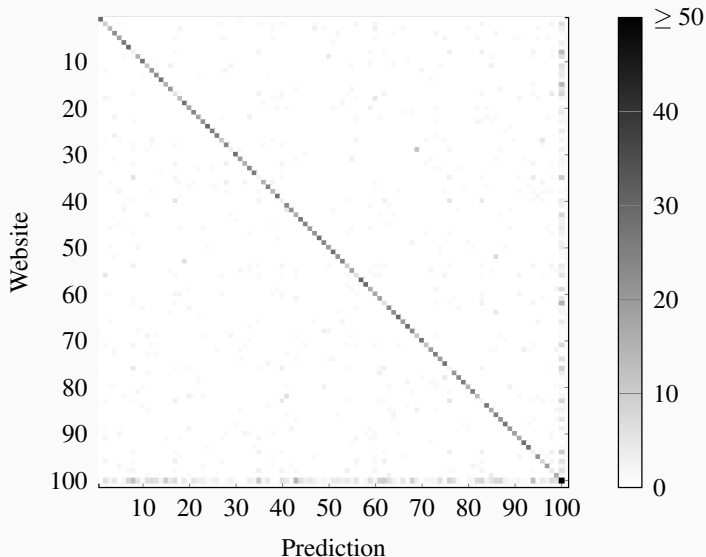
How large does the website have to be?



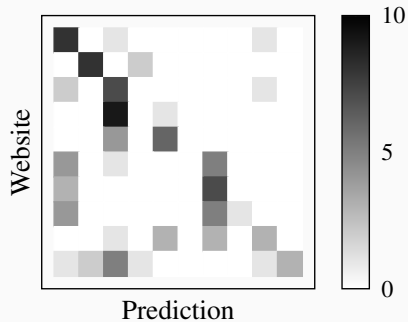
Video Fingerprinting on 10 different connections

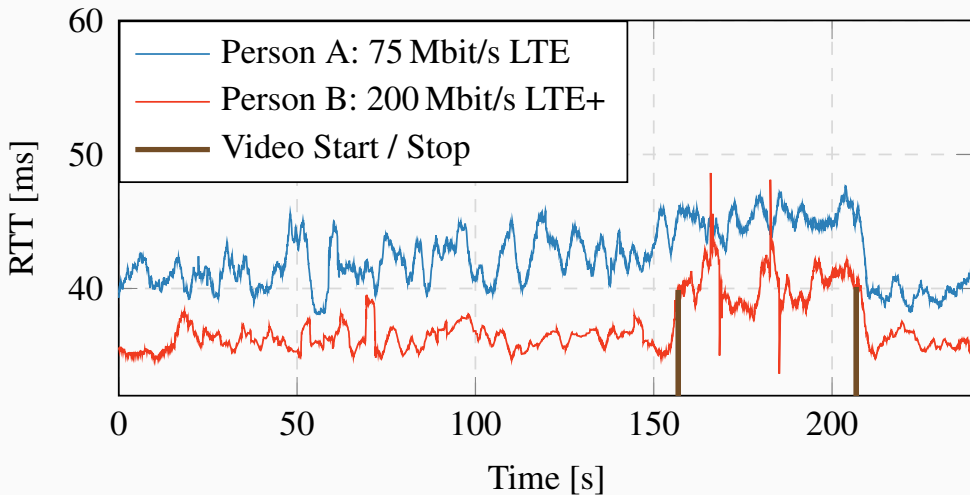


Top-100 Open-World Website Fingerprinting

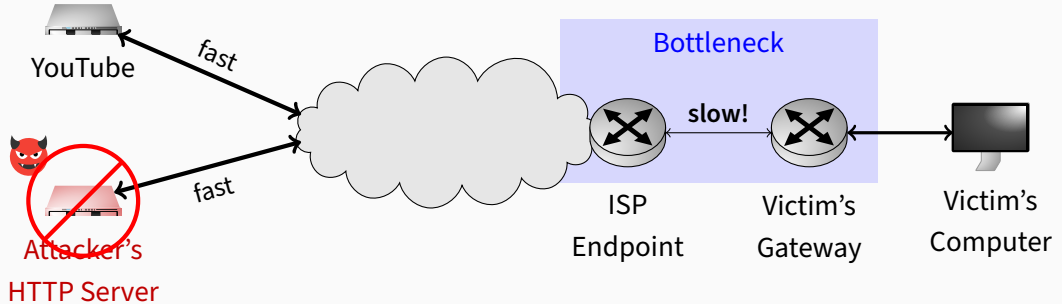


Cross-Connection Website Fingerprinting

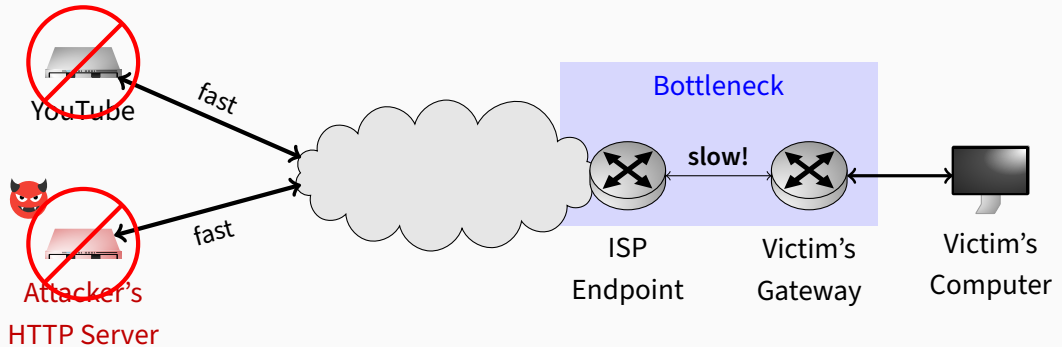




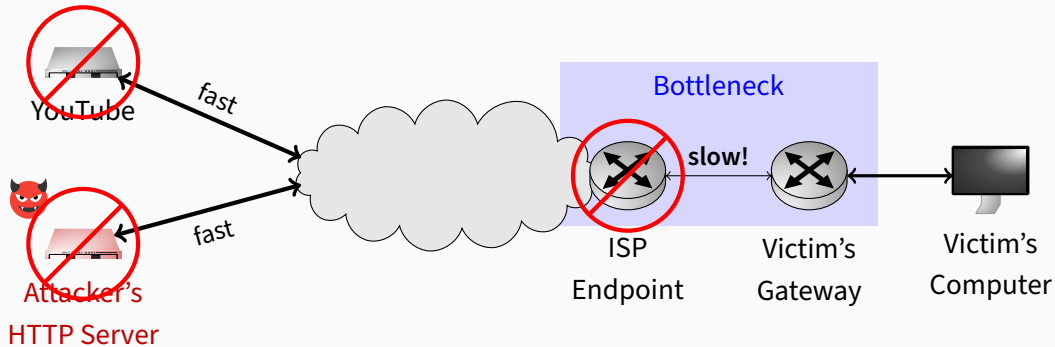
What about Mitigations?



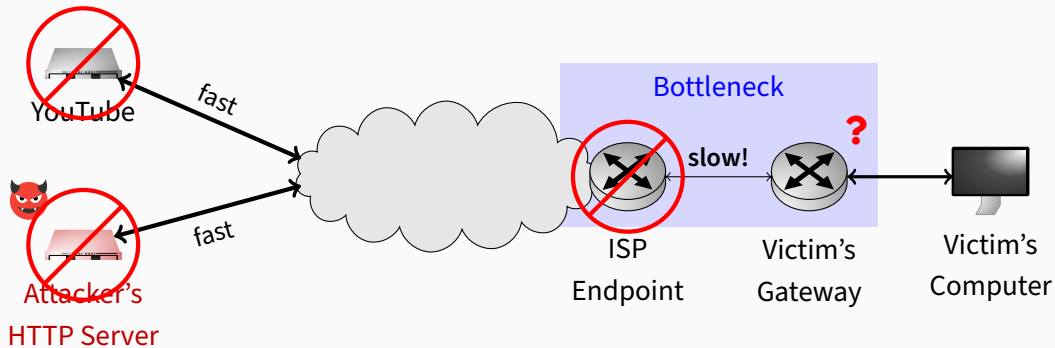
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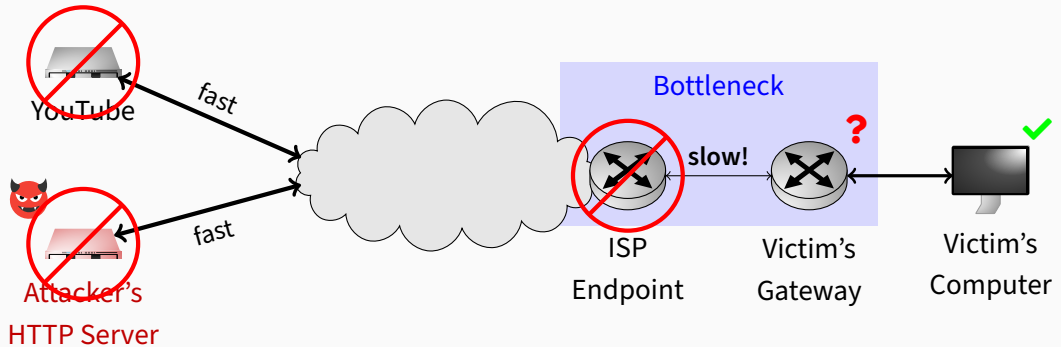
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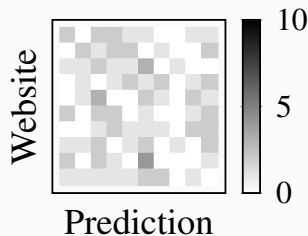
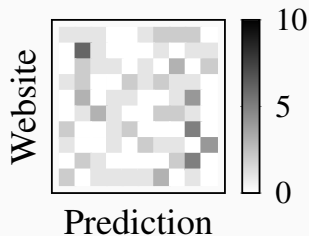
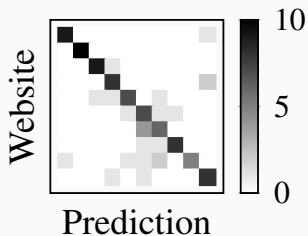
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What about Mitigations?



Impact of Noise on Website Fingerprinting



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- Many “remote” attacks can now be transformed to truly remote attacks
- We disclosed to Google / YouTube
 - they investigated the issue for several weeks
 - concluded that it is a generic problem

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- Paper + Demo: <https://snailload.com>

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- [AYR15] Daniel Arp, Fabian Yamaguchi, and Konrad Rieck. **Torben: A Practical Side-Channel Attack for Deanonymizing Tor Communication**. ASIA CCS. 2015.
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